

Computing Systems

Introduction to Operating Systems

Contents

- The History of Operating Systems
- Operating System Architecture
- Coordinating the Machine's Activities
- Handling Competition Among Processes
- Security

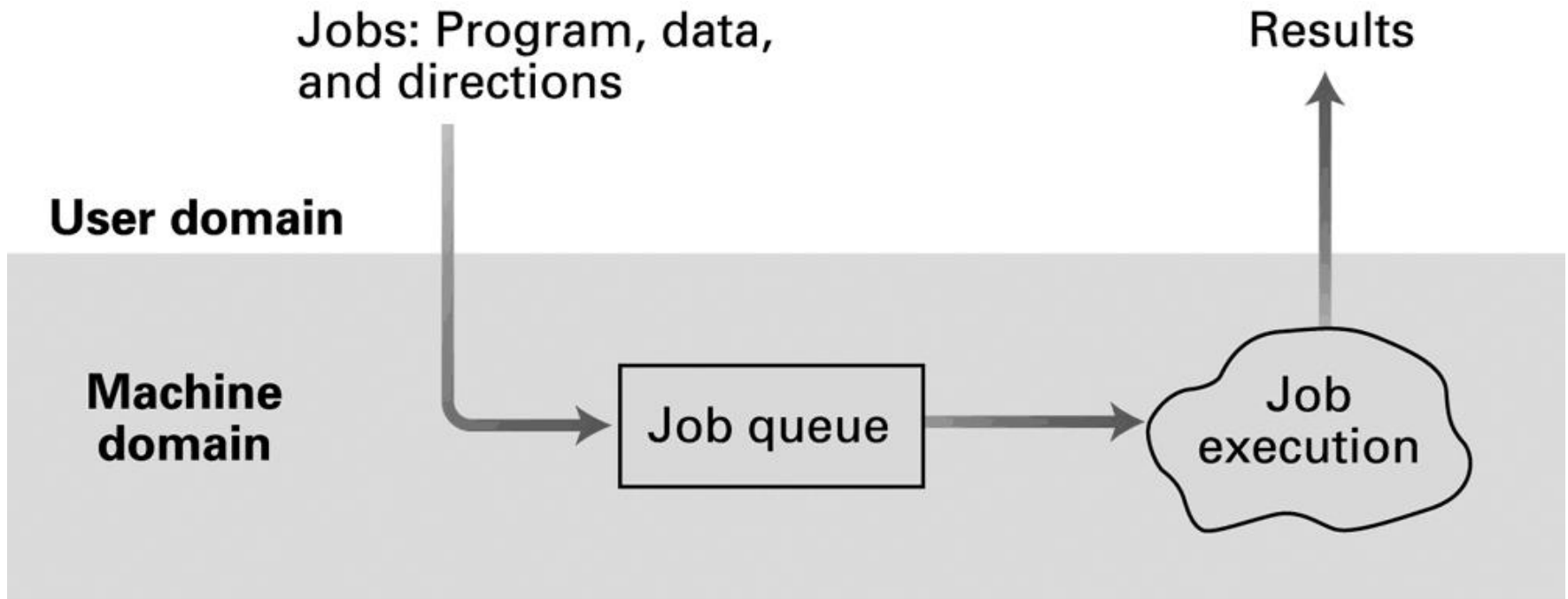
Functions of Operating Systems

- Oversee operation of computer
- Store and retrieve files
- Schedule programs for execution
- Coordinate the execution of programs

Evolution of Shared Computing

- Batch processing
- Interactive processing
 - Requires real-time processing
- Time-sharing/Multitasking
 - Implemented by Multiprogramming
- Multiprocessor machines

Batch processing



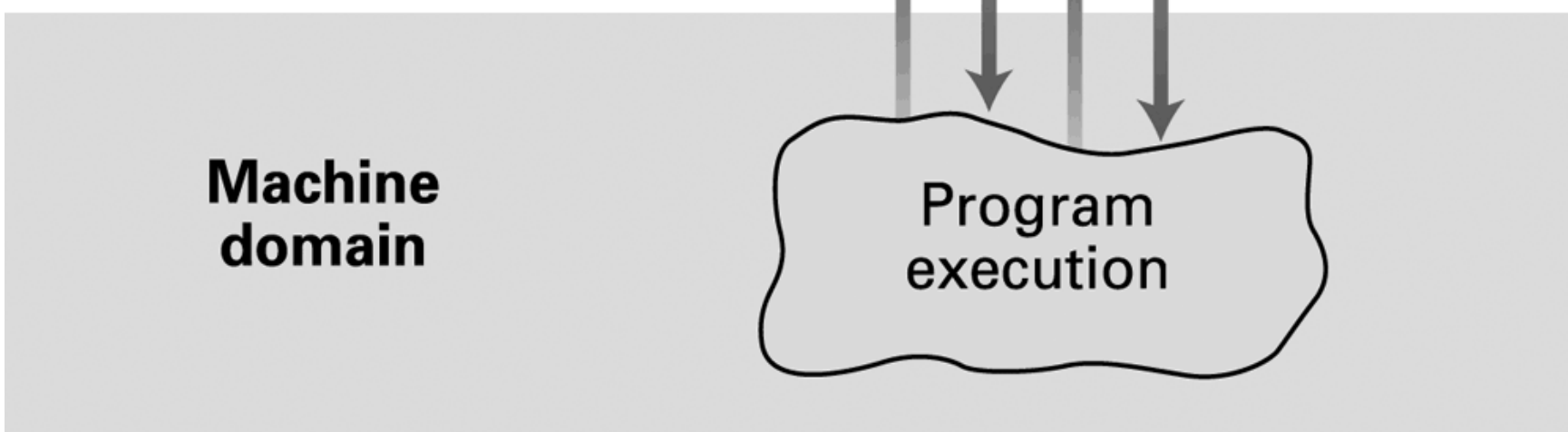
Interactive processing

Programs, data,
directions, and results

User domain

**Machine
domain**

Program
execution



The diagram illustrates the flow of information between the User domain and the Machine domain. The Machine domain is represented by a light gray rectangular area. Inside this area, there is a cloud-shaped box labeled 'Program execution'. Four vertical arrows connect the top of the 'Program execution' box to the text 'Programs, data, directions, and results' located above the Machine domain. The arrows alternate in direction: the first and third arrows point upwards, while the second and fourth arrows point downwards. The text 'User domain' is positioned to the left of the Machine domain, and the text 'Machine domain' is positioned to the left of the 'Program execution' box.

Time Sharing / Multitasking

- Users seeking services from same machine at the same time – time sharing
 - Implemented using a technique called multiprogramming (time is divided into multiple intervals, execution of one job is limited to a single time interval)
- Multiple terminals connected to same machine
 - Driven by the fact that in the past computers were very expensive
- When multiprogramming is applied to single-user environments is usually called multitasking

Multiprocessor Operating Systems

- Provide time sharing/multi-tasking capabilities by assigning different tasks to different processors as well as sharing the time of one single processor
- Problems to solve:
 - Load balancing – dynamically allocating tasks to the various processor so that all of them are used efficiently
 - Scaling – breaking tasks into sub-tasks compatible with the number of processors available
- Trend to develop a network wide operating system rather than networks of individual operating systems

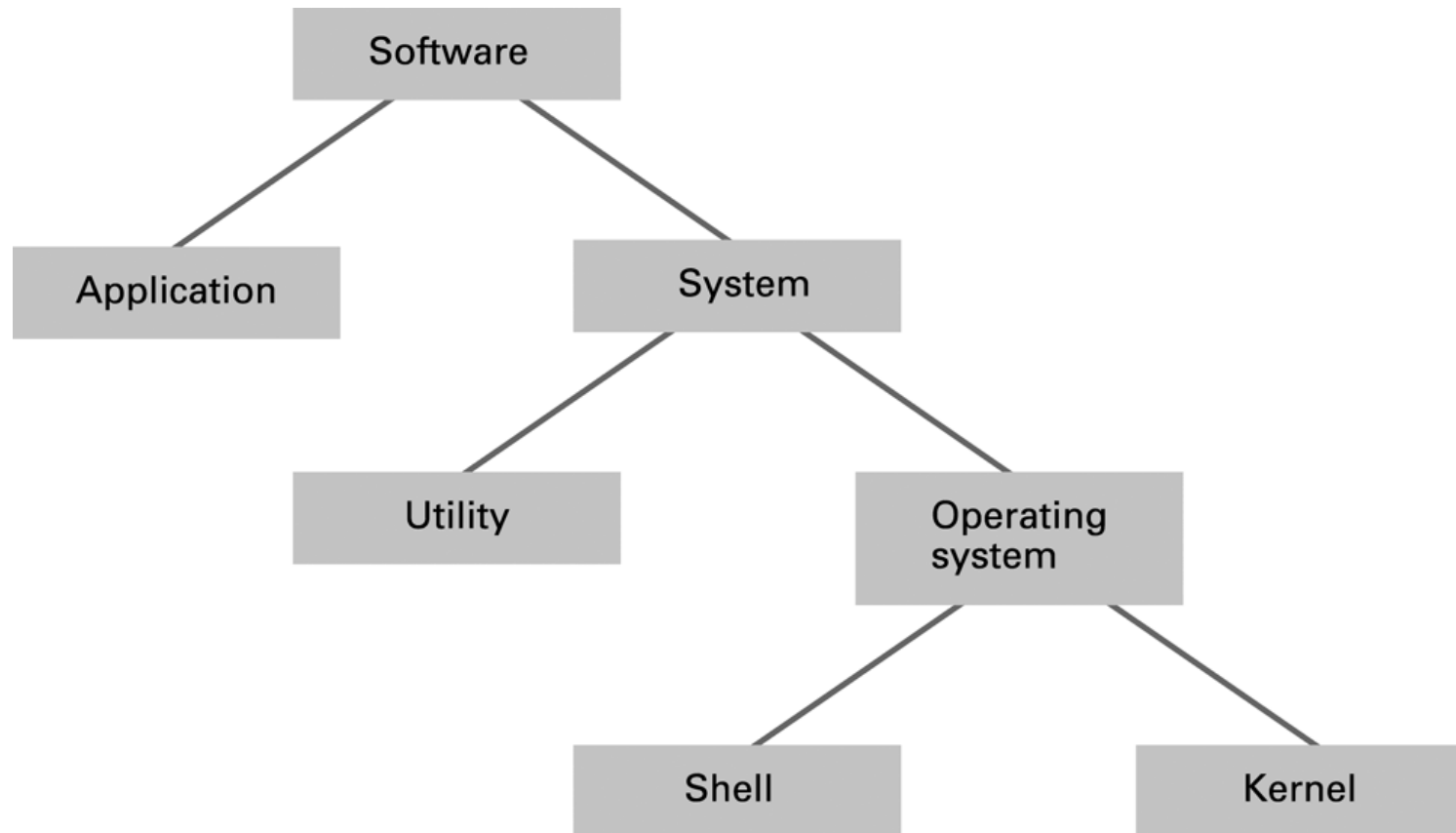
Embedded Operating Systems

- Used in hand held devices, mobile phones, cars, etc...
- Limited data storage and power conservation are the big challenges
- Examples: VxWorks, Windows CE (Pocket PC), Palm OS, Symbian, ThredX, RomDOS, etc...

Types of Software

- Application software
 - Performs specific tasks for users: spreadsheets, database systems, desktop publishing, program development, games, etc...
- System software
 - Provides infrastructure for application software
 - Consists of **operating system** and **utility software**

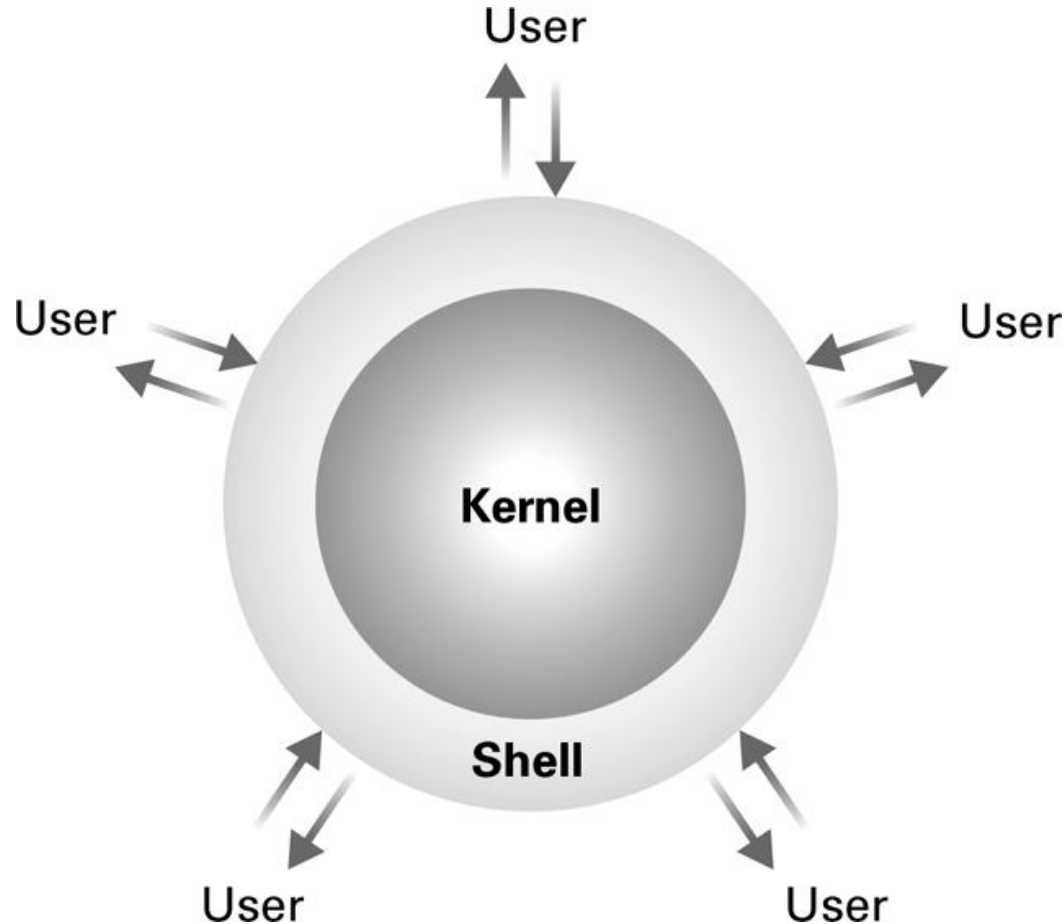
Software classification



Operating System Components

- **Shell:** Communicates with users
 - Text based
 - Graphical user interface (GUI)
- **Kernel:** Performs basic functions
 - File manager
 - Device drivers
 - Memory manager
 - Process manager (Scheduler, dispatcher, etc..)

The shell as an interface between users and the operating system



File Manager

- Role – coordinate the use of machine's mass storage facilities
- Hierarchical organization
 - **Directory** (or **Folder**): A user-created bundle of files and other directories (subdirectories)
 - **Directory Path**: A sequence of directories within directories
- Access/operations to files is provided by file manager via a **file descriptor**

Device Manager

- Part of OS presented as a collection of device drivers – specialized software that communicate with the controllers to carry out operations on peripheral devices connected to the computer
- Each driver is specifically designed for its type of device (e.g. printer, monitor, etc..) and translates generic requests into device specific sequence of operations

Memory Manager

- Has the task of coordinating the use of main memory – allocates/deallocates space in main memory
- When the total required memory space exceeds the physical available space.
 - May create the illusion that the machine has more memory than it actually does (**virtual memory**) by playing a “shell game” in which blocks of data (**pages**) are shifted back and forth between main memory and mass storage

Processes

- **Process:** The activity of executing a program
(NOT THE SAME THING AS A PROGRAM!!!)
 - Program – static set of directions (instructions)
 - Process – dynamic entity whose properties change as time progresses. It is an instance in execution of a program.
- **Process State:** Current status of the activity
 - Program counter
 - General purpose registers
 - Related portion of main memory

Process Manager

- **Scheduler** – the part of kernel in charge with the strategy for allocation/de-allocation of the CPU to each competing process
 - Maintains a record of all processes in the OS (via a **process table**), introduces new processes to this pool and removes the ones that completed
- **Dispatcher** is the component of the kernel that oversees the execution of the scheduled processes
 - Achieved by multiprogramming

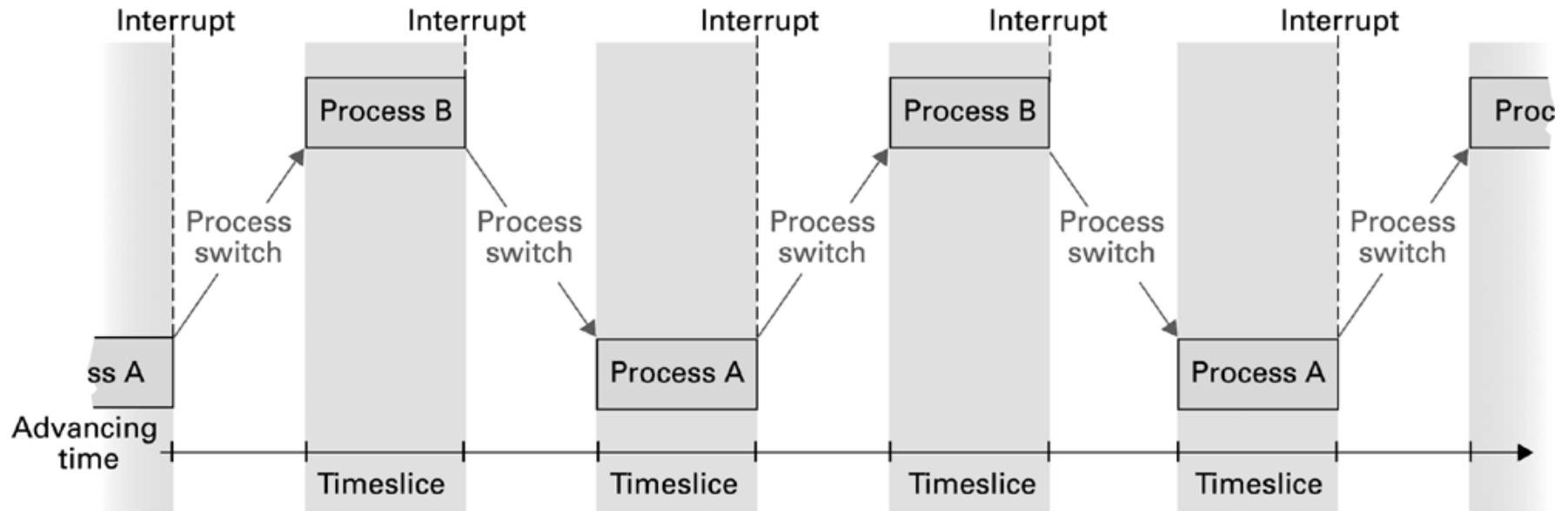
Scheduler

- **Scheduler:** Adds new processes to the process table and removes completed processes from the process table
- Process table contains
 - Memory area assigned to the process
 - Priority of the process
 - State of the process (ready or waiting)

Dispatcher

- **Dispatcher:** Controls the allocation of CPU (of time slices) to the processes in the process table
 - The end of a time slice is signaled by an interrupt.
 - Each process is allowed to execute for one time slice
- It performs “**process switch**” – procedure to change from one process to another
 - ProcessA → Dispatcher → ProcessB

Time-sharing between process A and process B



Handling Competition for Resources

Important task of OS is to allocate *resources* to the processes

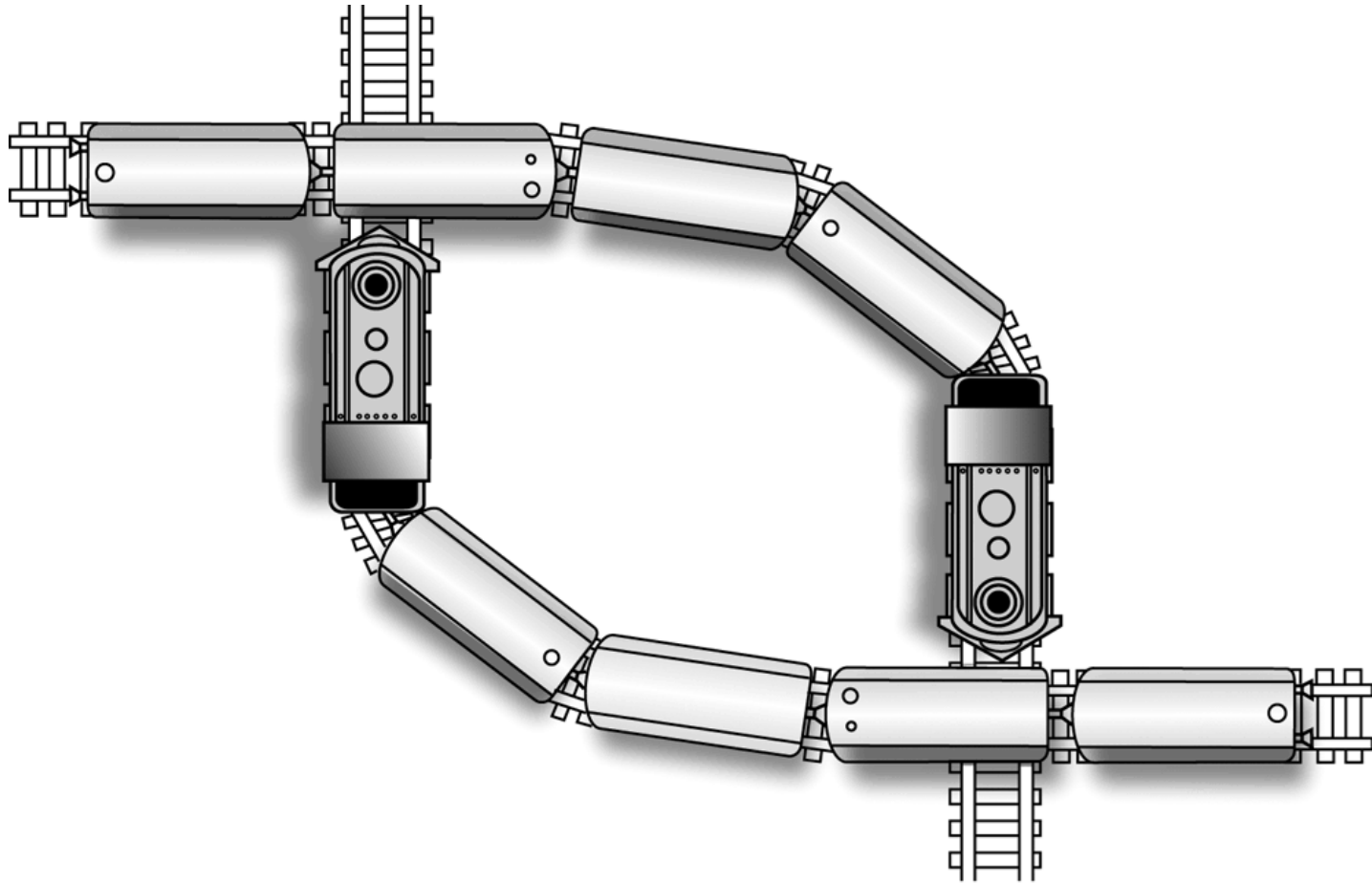
- **Semaphore:** A “control flag”
- **Mutual exclusion:** Requirement for proper implementation of a critical region so that only one process at a time will execute the sequence of instructions part of a critical region

Deadlock

Another problem of resource allocation :-

- Processes block each other from continuing
- Conditions required for deadlock
 1. Competition for non-sharable resources
 2. Resources requested on a partial basis
 3. An allocated resource can not be forcibly retrieved

A deadlock resulting from competition for nonshareable railroad intersections



Security

- One of the role of OS is to provide security
- Attacks from outside
 - Insecure passwords
 - Counter measures
 - Auditing software

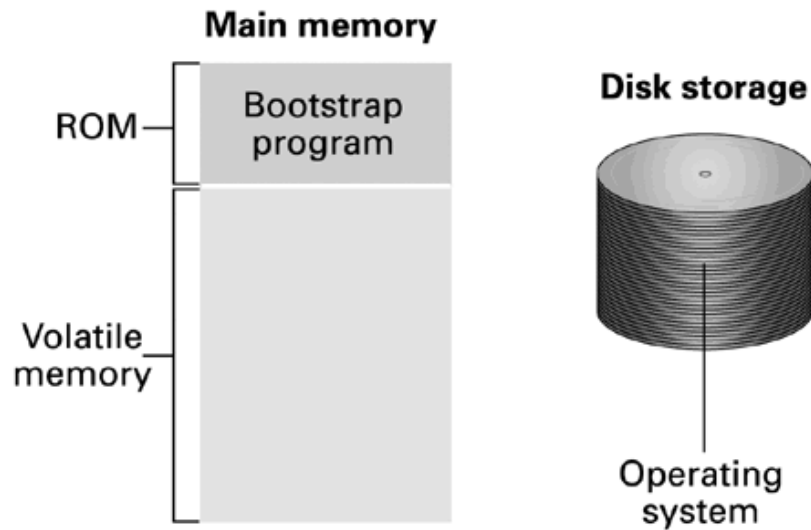
Security (continued)

- Attacks from within
 - Counter measures: Control process activities via privileged modes and privileged instructions
 - Examples on attacker SW:
 - Alters the timer of OS – extend its own time slice and dominate the machine
 - Access to peripheral devices directly – access to files that access would have been denied
 - Access memory cells outside its allowed area, it can read and alter data from other processes

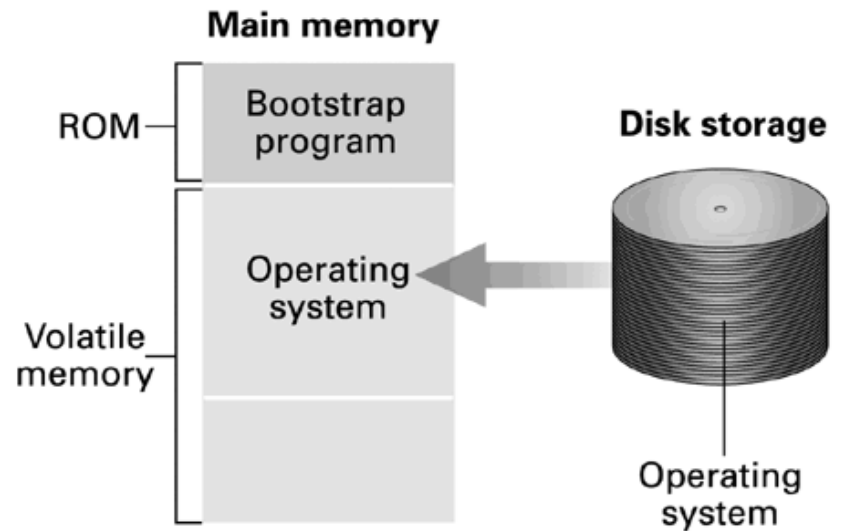
Getting OS Started (Bootstrapping)

- **Booting:** Procedure that transfers the OS from mass storage (permanent) into the main memory (volatile-thus empty when machine is turned on)
- **Bootstrap:** Program in ROM (example of firmware)
 - Run by the CPU when power is turned on (PC starts at pre-defined address when power is applied)
 - Transfers operating system from mass storage to main memory
 - Executes jump to operating system

The booting process



Step 1: Machine starts by executing the bootstrap program already in memory. Operating system is stored in mass storage.



Step 2: Bootstrap program directs the transfer of the operating system into main memory and then transfers control to it.